

**Solutions to Exercises in
Game Theory
Chapter 6**

1. [Typo: The formula for the optimal bid should be

$$\beta(x) = \frac{1}{G(x)} \int_0^x yg(y) dy$$

with y inside the integral and x as upper bound for the integral.] To show that β is a symmetric equilibrium strategy, we let the bidder act as if the value was z (rather than x). The expected payoff is then

$$G(z)(x - \beta(z))$$

(the probability of $\beta(z)$ being largest or (assuming monotonicity of β) of z being largest value, multiplied with the gain if winning), and this can be rewritten as

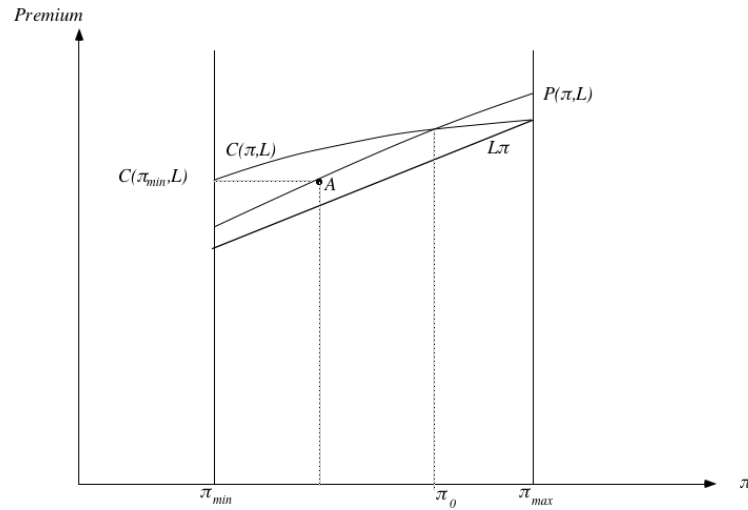
$$\begin{aligned} G(z)(x - \beta(z)) &= G(z)x - \int_0^z yg(y) dy \\ &= G(z)x - G(z)z + \int_0^z G(y) dy = G(z)(x - z) + \int_0^z G(y) dy, \end{aligned}$$

where we have used integration by parts. Subtracting this from the expected gain using x , which is $\int_0^x yg(y) dy$, we get the difference

$$G(z)(z - x) - \int_x^z G(y) dy,$$

which is easily seen to be ≥ 0 for all values of z .

2. The figure below contains three curves, namely for each p (in the text p) (1) the average loss of an individual with risk p , (2) the average loss $C(p, L)$ of an individual with risk $\geq p$, and (3) the willingness to pay $P(p, L)$ of an individual with risk p .



The individuals for which $C(p; L) > P(p, L)$ will not buy insurance since they consider it as too expensive.

3. In the first step of this procedure (where we determine $t(a)$ for each e), we solve the problem of maximizing

$$\sum_{h=1}^r p_h(e)(y_h - t_h)$$

over all (t_1, \dots, t_r) such that

$$\begin{aligned} \sum_{h=1}^r p_h(e)(v(t_h) - w(e)) &= 0, \\ \sum_{h=1}^r p_h(e)(v_h(t_h) - w(e)) &\geq \sum_{h=1}^r p_h(e')(v(t_h) - w(e')), \text{ all } e'. \end{aligned}$$

where the first constraint a participation condition (the 0 on the right-hand side represents the expected value of alternative engagement) and the second is the incentive compatibility constraint (the agent must be induced to deliver the effort e).

Let $t^* = (t_1^*, \dots, t_r^*)$ be a solution, and let $K(e)$ be set of e' such that the last condition is fulfilled with equality. Restricting to the corresponding equations gives a maximization problem with first order conditions

$$p_h(e) + \lambda v'(t_h^*) + \sum_{e' \in K(e)} \mu(e') v'(t_h^*) (p_h(e) - p_h(e')) = 0,$$

or

$$\frac{1}{v'(t_h^*)} = -\lambda - \sum_{e' \in K(e)} \mu(e') \frac{p_h(e) - p_h(e')}{p_h(e)},$$

for $h = 1, \dots, r$, where λ and $\mu(e')$, $e' \in K(a)$, are Lagrangian multipliers.

4. In a symmetric Bayesian Nash equilibrium, the bid function $b(x)$ (depending on the signal x received by the individual), assumed to be monotone, is the same for both, and the payoff is

$$\pi(b, x) = \int_0^{\beta^{-1}(b)} (v(x, y) - \beta(y))g(y|x) dy$$

where $g(y|x)$ is the conditional density of the signal received by the other individual, given that x is the highest signal – the first individual wins if her bid is the highest and pays only the bid of the other individual. Writing the integral as

$$\int_0^{\beta^{-1}(b)} (v(x, y) - v(y, y))g(y|x) dy$$

(using the both have the same bidding function), we use that $v(x, y)$ is $>$ or $<$ than $v(y, y)$ depending on whether $x > y$ or $x < y$. To maximize π , we need to keep all the positive and discard the negative contributions to the integral, and this is obtained by choosing $\beta^{-1}(b) = x$ or $b = \beta(x)$.

Suppose that

$$v(x, y) = \frac{1}{3}x + \frac{2}{3}y$$

(valuation depends with weight 1/3 on own signal and with weight 2/3 on that of the other bidder). Then the valuation of bidder 1 is greater than that of bidder 2 if and only if $y > x$. In this case bidder 2 has the highest bid and wins the auction, but bidder 2 values it higher, meaning that the second-price auction is not efficient.

Let $x > y$ be arbitrary, and consider the path from (y, x) to (x, y) given by

$$\gamma(t) = (1 - t)(y, x) + t(x, y), \quad t \in [0, 1]$$

By the mean value theorem, there is some t^0 such that

$$\frac{dv(\gamma(t))}{dt}(t^0) = v(x, y) - v(y, x)$$

The left-hand side can be written as

$$v'_1(\gamma(t^0))(x - y) + v'_2(\gamma(t^0))(y - x) = (v'_1 - v'_2)(x - y),$$

and by single-crossing and our assumption $x > y$, we get that $v(x, y) > v(y, x)$, so that highest valuation also has highest bid.

5. If the buyer must propose a price initially, then utility is $\tau_B q - p$ if $p \geq \tau_S q$ and 0 otherwise, so that trade will occur whenever $\tau_B \geq \tau_S$, and the proposed price is $\tau_S \frac{Q}{2}$.

When the seller proposes a price initially, the buyer will accept when expected value of q exceeds the proposed price, and since the quality must be in the interval $\left[0, \frac{p}{\tau_S}\right]$ with mean $\frac{p}{2\tau_S}$ and the buyer will accept when

$$p \leq \tau_B \frac{p}{2\tau_S}$$

or $\tau_B \geq 2\tau_S$. In this case, the seller will obtain maximal payoff if the price is set as $\tau_S Q$.